

Robust Procurement Design

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Motivation

- Procurement

central economic problem

(a) public

(b) private

- Difficulty in designing contracts

Provider typically **privately informed** about cost

Motivation

- Standard approach: Bayesian
 - Buyer has conjecture/belief over
 - (a) supplier's cost
 - (b) gross value for procuring output
 - maximizes under conjecture
- CS (robust) approach
 - buyer has no conjecture
 - worst-case optimality

- **Our Approach (alternative form of robustness)**
 - buyer has conjecture but does not fully trust it
 - prepares for the worst (in case conjecture is wrong)
 - uses conjecture to select optimal mechanism among worst-case optimal ones
 - Lexicographic approach
 - (a) political/hierarchical constraints
 - (b) attitude towards ambiguity

This Paper: Results

- Robust design of procurement contracts
- Uncertainty only over cost:
 - Baron Myerson with **quantity floor**
- Uncertainty over both cost and demand
 - **upward quantity adjustment for high cost**
 - **downward quantity adjustment for low/intermediate costs**
- Robustly optimal mechanism sensitive only to
 - conjecture over demand and cost
 - lowest admissible demand
- CS wrt buyer's uncertainty and policy recommendations

Literature – Incomplete

- Baron-Myerson, Ecma (1982), Lewis and Sappington (1988), Armstrong, JET (1999)
 - **Bayesian**
- Bergemann-Heumann-Morris, wp (2024)
 - **different robustness criterion (competitive ratio)**
- Guo-Shmaya, wp (2024)
 - **different robustness criterion (min-max regret)**
- Garrett, GEB (2014)
 - **different robustness criterion (worst-case optimality)**
- Mishra-Patel, wp (2024)
 - **undominated mechanisms**
- Dworzak-Pavan, Ecma (2023)
 - **robust information design**

Plan

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Model

- **Players**

- Buyer (principal/government)
- Seller (agent/monopolist)

- **Choices**

- output $q \geq 0$
- transfer $t \geq 0$

- **Payoffs**

- Social value of q :

$$V^*(q) = \int_0^q P^*(s) ds$$

- Cost: θq

θ drawn from ab. cont. F^* with $f^*(\theta) > 0$ over $\Theta = [\underline{\theta}, \bar{\theta}]$

- Ex-post welfare ($\alpha \in [0, 1]$)

$$V^*(q) - t + \alpha(t - \theta q)$$

- Ex-ante welfare:

$$\int (V^*(q) - t + \alpha(t - \theta q)) dF^*(\theta)$$

- **Asymmetric information**

- θ : monopolist's private information

● Uncertainty/Robustness

- government not sure about conjecture (V^*, F^*)
- concerned demand and cost may be $(V, F) \neq (V^*, F^*)$

● Admissible sets

- \mathcal{V} : set of possible consumer (gross) value fns
 - each $V \in \mathcal{V}$ strictly increasing, strictly concave, differentiable
- \mathcal{P} : set of corresponding *inverse demand functions*
- $P^* \in \mathcal{P}$ and $V^* \in \mathcal{V}$
- lowest (inverse) demand function: \underline{P}

- any $q \geq 0$ and $P \in \mathcal{P}$

$$P(q) \geq \underline{P}(q)$$

- \underline{P} : strictly decreasing and continuous and s.t.

$$\lim_{q \rightarrow 0^+} \underline{P}(q) > \bar{\theta}$$

Model

- **(Direct) mechanism** $M = (q, t)$
 - quantity schedule $q : \Theta \rightarrow \mathbb{R}_+$
 - (total) transfer schedule $t : \Theta \rightarrow \mathbb{R}$
- $M = (q, t)$ IC and IR iff, for all θ, θ' ,

$$t(\theta) - \theta q(\theta) \geq t(\theta') - \theta q(\theta')$$

with

$$t(\theta) - \theta q(\theta) \geq 0$$

- \mathcal{M} : set of IC and IR mechanisms

- Given IC and IR mechanism $M = (q, t)$, **ex-ante welfare** under (V, F)

$$W(M; V, F) := \int w(\theta, M; V) F(d\theta)$$

where **ex-post welfare**:

$$w(\theta, M; V) := V(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)$$

total surplus:

$$V(q(\theta)) - \theta q(\theta)$$

“**rent**” to type θ :

$$u(\theta) := t(\theta) - \theta q(\theta)$$

Definition

Given any IC and IR $M = (q, v)$, **welfare guarantee**

$$G(M) := \inf_{V \in \mathcal{V}, F \in \mathcal{F}} W(M; V, F)$$

- **Short-list:**

$$\mathcal{M}^{\text{SL}} := \{M \in \mathcal{M} : G(M) \geq G(M') \forall M' \in \mathcal{M}\}$$

(set of IC and IR mechanisms for which guarantee is maximal)

- **Government's problem**

- choose mechanism from \mathcal{M}^{SL} maximizing welfare under conjecture (V^*, F^*)

Definition

Mechanism $M \in \mathcal{M}^{\text{SL}}$ **robustly optimal** iff, for every $M' \in \mathcal{M}^{\text{SL}}$,

$$W(M; V^*, F^*) \geq W(M'; V^*, F^*)$$

- Robustly optimal mechanisms maximize ex-ante welfare **under conjecture (V^*, F^*)** over all worst-case-optimal mechanisms

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Baron Myerson

Baron Myerson

- $\partial u(\theta)$: set of all subgradients of u at θ
 - $x \in \mathbb{R}$ is subgradient iff for all θ' , $u(\theta') \geq u(\theta) + x(\theta' - \theta)$

Lemma

Mechanism $M = (q, u)$ IC and IR iff u is convex, non-increasing, and s.t, for all $\theta \in \Theta$, $-q(\theta) \in \partial u(\theta)$, with $u(\bar{\theta}) \geq 0$. Equivalently,

- 1 q non-increasing
- 2 for all $\theta \in \Theta$,

$$u(\theta) = u(\bar{\theta}) + \int_{\theta}^{\bar{\theta}} q(z) dz$$

- 3 $u(\bar{\theta}) \geq 0$

- θ 's **virtual cost** (under conjecture F^*):

$$z^*(\theta) := \theta + (1 - \alpha) \frac{F^*(\theta)}{f^*(\theta)}$$

Proposition

Assume z^* non-decreasing. (Bayesian) optimal mechanism $M^{BM} = (q^{BM}, u^{BM})$ s.t., for all θ ,

$$q^{BM}(\theta) := \arg \max_q \{V^*(q) - z^*(\theta)q\}$$

$$u^{BM}(\theta) = \int_{\theta}^{\bar{\theta}} q^{BM}(z) dz$$

- **Proof:** Familiar IC analysis \rightarrow ex-ante welfare under conjecture (V^*, F^*)

$$\int_{\underline{\theta}}^{\bar{\theta}} [V^*(q(\theta)) - z^*(\theta)q(\theta)] dF^*(\theta)$$

- $q^{BM}(\theta)$ maximizes virtual surplus $V^*(q) - z^*(\theta)q$ point-wise
- z^* non-decreasing $\Rightarrow q^{BM}$ non-increasing $\Rightarrow (q^{BM}, u^{BM})$ IC and IR

- FB-efficiency (under conjecture (V^*, F^*)):

$$P^*(q^{FB}(\theta)) = \theta$$

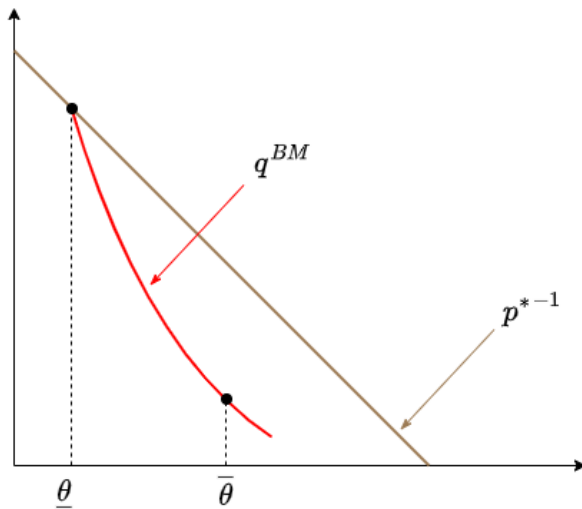
- BM schedule (second-best efficiency)

$$P^*(q^{BM}(\theta)) = z^*(\theta) := \theta + (1 - \alpha) \frac{F^*(\theta)}{f^*(\theta)}$$

- Hence,

- **no distortion “at top”**, i.e., for most efficient type, $\underline{\theta}$
- **downward distortions for all $\theta > \underline{\theta}$**

Baron Myerson



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Short List

Definition

Mechanism $M = (q, u)$ **worst-case optimal** iff, for any IC and IR $M' = (q', u') \in \mathcal{M}$,

$$G(M') := \inf_{V \in \mathcal{V}, F \in \mathcal{F}} W(M'; V, F) \leq \inf_{V \in \mathcal{V}, F \in \mathcal{F}} W(M; V, F) := G(M)$$

Worst-case optimality

- Let

$$q_\ell := \arg \max_q \{ \underline{V}(q) - \bar{\theta}q \}$$

denote efficient output when $V = \underline{V}$ and $\theta = \bar{\theta}$

Lemma

For any IC and IR mechanism $M = (q, u) \in \mathcal{M}$

$$G(M) = \inf_{\theta \in \Theta} \{ \underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta) \}$$

and

$$G(M) \leq G^* := \underline{V}(q_\ell) - \bar{\theta}q_\ell.$$

Worst-case optimality

Proof: For any IC and IR mechanism $M = (q, u) \in \mathcal{M}$

$$\begin{aligned}W(M; V, F) &:= \int \{V(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)\} F(d\theta) \\ &\geq \int \{\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)\} F(d\theta) \\ &\geq \inf_{\theta \in \Theta} [\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)]\end{aligned}$$

Hence,

$$G(M) \geq \inf_{\theta \in \Theta} [\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)]$$

Because $\underline{V} \in \mathcal{V}$ and, for each θ , Dirac distribution at θ is in \mathcal{F}

$$G(M) \leq \inf_{\theta \in \Theta} [\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)]$$

Hence,

$$G(M) = \inf_{\theta \in \Theta} [\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)]$$

Because $u(\bar{\theta}) \geq 0$ and

$$\underline{V}(q(\bar{\theta})) - \bar{\theta}q(\bar{\theta}) \leq \underline{V}(q_\ell) - \bar{\theta}q_\ell := G^*$$

$$G(M) \leq G^*$$

Short List: Characterization

- Let

$$\mathcal{M}^{\text{SL}} := \{M \in \mathcal{M} : G(M) \geq G(M') \forall M' \in \mathcal{M}\}$$

Lemma

$M = (q, u) \in \mathcal{M}^{\text{SL}}$ iff q non-increasing and, for any θ

$$u(\theta) = \int_{\theta}^{\bar{\theta}} q(y) dy$$

and

$$\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q(y) dy \geq G^*$$

Short List: Characterization

Proof:

- $M = (q, u)$ IC and IR:
 - (a) q nondecreasing
 - (b) $u(\theta) = u(\bar{\theta}) + \int_{\theta}^{\bar{\theta}} q(y)dy$, with $u(\bar{\theta}) \geq 0$
- ONLY IF:
 - $M_L = (q_L, u_L)$ w. $q_L(\theta) = q_\ell$ and $t_L(\theta) = \bar{\theta}q_\ell$ all θ
 - (i) IC and IR
 - (ii) $u_L(\theta) = (\bar{\theta} - \theta)q_\ell$
 - (iii) $w(\theta, M_L; \underline{V}) := \underline{V}(q_\ell) - \theta q_\ell - (1 - \alpha)u_L(\theta) = G^* + \alpha(\bar{\theta} - \theta)q_\ell$
 - (iv) $G(M_L) = G^*$
 - $M \in \mathcal{M}^{\text{SL}}$ **only if** $G(M) = G^*$
 - Because $w(\theta, M; \underline{V}) := \underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha)u(\theta)$
 - (1) $u(\bar{\theta}) = 0$
 - (2) $\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q(y)dy \geq G^*$ all θ
- IF PART: immediate

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Robust Optimality

Robust Optimality: Full Program

- Recall def of virtual cost under conjecture (F^*, V^*)

$$z^*(\theta) := \theta + (1 - \alpha) \frac{F^*(\theta)}{f^*(\theta)}$$

- Robustly optimal schedule q^{OPT} solves

$$\max_q \int_{\underline{\theta}}^{\bar{\theta}} [V^*(q(\theta)) - z^*(\theta)q(\theta)] dF^*(\theta)$$

s.t.

q non-increasing

$$\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q(y) dy \geq G^* \quad \forall \theta \in \Theta$$

Robust Optimality: Relaxed Program

- Relaxation:

$$\max_{\substack{q \\ \underline{\theta}}} \int_{\underline{\theta}}^{\bar{\theta}} \left[V^*(q(\theta)) - z^*(\theta)q(\theta) \right] dF^*(\theta)$$

s.t.

q non-increasing

$$q(\theta) \geq q_\ell \quad \forall \theta \in \Theta$$

$$q(\bar{\theta}) = q_\ell$$

Robust Optimality: Relaxed Program

- **Proof:** $M = (q, u) \in \mathcal{M}^{\text{SL}}$ only if, for all θ ,

$$\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q(y) dy \geq G^*$$

- Because $G^* = \max_q \{ \underline{V}(q) - \bar{\theta}q \} = \underline{V}(q_\ell) - \bar{\theta}q_\ell$
- $\Rightarrow q(\bar{\theta}) = q_\ell$
- q non-increasing $\Rightarrow q(\theta) \geq q_\ell$ all θ

Robust Optimality: BM with Floor

- Let

$$q^{BM}(\theta) := \arg \max_q \{V^*(q) - z^*(\theta)q\}$$

$$q^*(\theta) := \begin{cases} \max\{q^{BM}(\theta), q_\ell\} & \text{if } \theta \neq \bar{\theta} \\ q_\ell & \text{if } \theta = \bar{\theta} \end{cases}$$

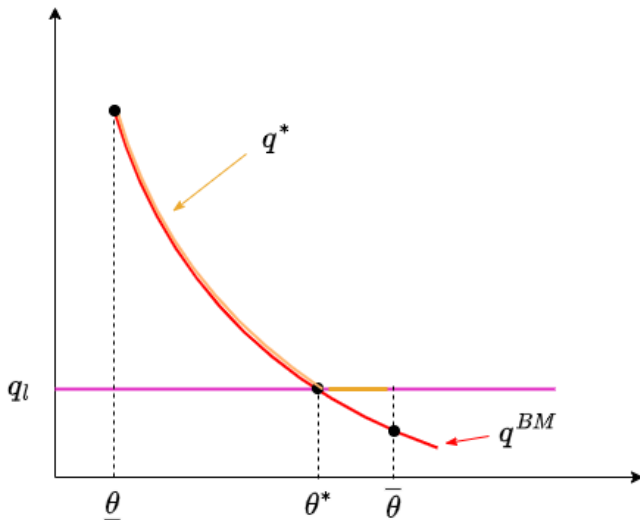
and

$$u^*(\theta) = \int_{\theta}^{\bar{\theta}} q^*(y) dy$$

Proposition

Suppose z^* non-decreasing and $V^* = \underline{V}$. Mechanism $M^* = (q^*, u^*)$ robustly optimal

Robust Optimality: BM with Floor



Robust Optimality: BM with Floor

Proof

- $M^* = (q^*, u^*)$ solves relaxed program
 - $q^*(\theta)$ maximizes $V^*(q) - z^*(\theta)q$ under constraint $q \geq q_\ell$
 - z^* non-decreasing $\Rightarrow M^*$ IC and IR
- Ex-post welfare under M^* and \underline{V}

$$w(\theta; M^*, \underline{V}) := \underline{V}(q^*(\theta)) - \theta q^*(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q^*(y) dy$$

- $V^* = \underline{V} \Rightarrow w(\cdot; M^*, \underline{V})$ non-increasing, with $w(\bar{\theta}; M^*, \underline{V}) = \underline{V}(q_\ell) - \bar{\theta}q_\ell := G^*$
- Robustness constraints $w(\theta; M^*, \underline{V}) \geq G^*$ satisfied
- $\Rightarrow M^* = (q^*, u^*)$ **robustly optimal**

Robust Optimality: BM with Floor

- When demand known (only uncertainty over cost)
 - robust mechanism is Baron **Myerson with floor**
 - **efficiency at both bottom and top**
 - possibility that cost is less favorable than conjectured → **more output**
 - flat mechanisms never optimal:

$$\theta^* := \inf \{ \theta \in \Theta : q^*(\theta) = q_\ell \} > \underline{\theta}$$

- contrast w. standard “CS” approach to robustness ($q(\theta) = q_\ell$ all θ)

Robust Optimality: general case

- For any $M = (q, u)$ s.t. $u(\theta) = \int_{\theta}^{\bar{\theta}} q(y) dy$, ex-post welfare under lowest demand

$$\underline{W}(\theta, q) := \underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q(y) dy$$

- Robustness:

$$\underline{W}(\theta, q) \geq G^* \quad \forall \theta \in \Theta$$

Robust Optimality: general case

Theorem

Suppose z^* non-decreasing

(1) BM-floor mechanism $M^* \equiv (q^*, u^*)$ robustly optimal iff

$$\theta_m := \sup\{\theta : \theta \in \arg \inf_{\theta'} \underline{W}(\theta', q^*)\} = \bar{\theta}$$

(2) If $\theta_m < \bar{\theta}$, then $\theta_m < \theta^* := \inf\{\theta \in \Theta : q^*(\theta) = q_\ell\}$ and

(a) $q^{\text{OPT}}(\theta) = q_\ell$ for all $\theta \in [\theta^*, \bar{\theta}]$

(b) $q^{\text{OPT}}(\theta) \leq q^{\text{BM}}(\theta)$ for almost all $\theta \leq \theta^*$

(inequality strict over positive-measure $I \subseteq [\underline{\theta}, \theta^*]$)

(3) if $\theta_m \in (\underline{\theta}, \bar{\theta})$, $\alpha = 0$, and z^* increasing, $q^{\text{OPT}}(\theta) = q^{\text{BM}}(\theta)$ for $\theta \in (\underline{\theta}, \theta_m)$

Robust Optimality: general case

- (Qualitative) properties driven by

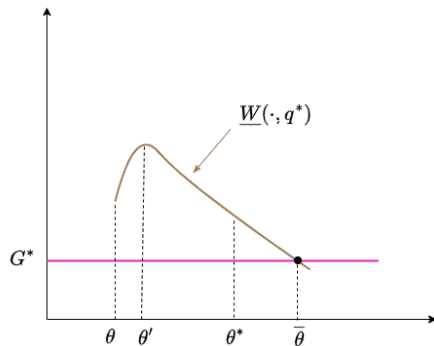
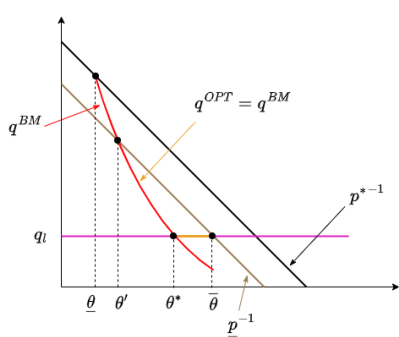
$$\underline{W}(\theta, q^*) := \underline{V}(q^*(\theta)) - \theta q^*(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q^*(y) dy$$

Lemma (Monotonicity)

For any positive, non-decreasing q

- 1 When $q(\theta) \leq \underline{P}^{-1}(\theta)$ all $\theta \in I \subset [\underline{\theta}, \theta^*]$, $\underline{W}(\cdot, q)$ non-increasing
(decreasing if $\alpha > 0$, or q decreasing and $q(\theta) < \underline{P}^{-1}(\theta)$ all $\theta \in I$)
- 2 When $q(\theta) > \underline{P}^{-1}(\theta)$ all $\theta \in I \subset [\underline{\theta}, \theta^*]$, and $\alpha = 0$, $\underline{W}(\cdot, q)$ non-decreasing
(increasing if q decreasing)

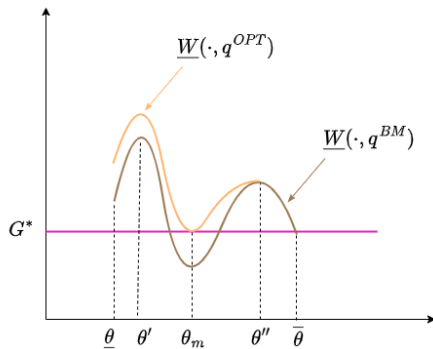
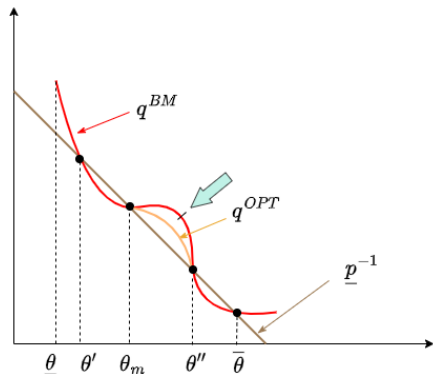
Robust Optimality: BM-floor optimal



Robust Optimality: BM-floor optimal

- When $\alpha = 0$, $M^* \equiv (q^*, u^*)$ robustly optimal if $\underline{W}(\underline{\theta}, q^*) \geq G^*$
- Equivalently, when $\underline{V}(q^*(\underline{\theta})) - \underline{\theta}q^*(\underline{\theta}) - \int_{\underline{\theta}}^{\bar{\theta}} q^*(y) dy \geq \underline{V}(q_\ell) - \bar{\theta}q_\ell$
- $q^*(\underline{\theta}) = \arg \max_q \{V^*(q) - \underline{\theta}q\}$
- Hence inequality holds when $\|V^*, \underline{V}\|$ small

Robust Optimality: BM with floor and intermediate downward adjustments



Optimality of “plateau” for high costs

Lemma

Suppose z^* non-decreasing and $\theta_m \neq \bar{\theta}$. Then $q^{\text{OPT}}(\theta) = q_\ell$ for all $\theta \in [\theta^*, \bar{\theta}]$

- Minimizing quantity over $[\theta^*, \bar{\theta}]$
 - increases $\underline{W}(\theta, q)$ for all $\theta \leq \theta^*$ (by reducing rents)
 - possibly decreases $\underline{W}(\theta, q)$ for $\theta > \theta^*$ (by reducing $\underline{V}(q) - \theta q$)
 - However, $\underline{W}(\cdot, q^*)$ decreasing over $[\theta^*, \bar{\theta}]$ w. $\underline{W}(\bar{\theta}, q^*) = G^*$
- For $\theta \in [\theta^*, \bar{\theta}]$
 - q_ℓ maximizes $V^*(q) - z^*(\theta)q$ over $[q_\ell, +\infty)$
- Hence under any $M^{\text{OPT}} = (q^{\text{OPT}}, u^{\text{OPT}})$

$$q^{\text{OPT}}(\theta) = q_\ell \quad \forall \theta \in [\theta^*, \bar{\theta}]$$

Suboptimality of upward adjustments for low costs

Lemma

Suppose z^* non-decreasing and $\theta_m \neq \bar{\theta}$. Then $q^{\text{OPT}}(\theta) \leq q^{\text{BM}}(\theta)$ for almost all $\theta \in [\underline{\theta}, \theta^*)$, with inequality strict over positive-measure $I \subseteq [\underline{\theta}, \theta^*)$

Sub-optimality of upward adjustments for low costs

- Take any $M = (q, u) \in \mathcal{M}^{\text{SL}}$ s.t.
 - $q(\theta) > q^*(\theta) = q^{\text{BM}}(\theta)$ over positive-measure $I \subseteq [\underline{\theta}, \theta^*)$
 - $q(\theta) = q_\ell$ for all $\theta \geq \theta^*$
- Take $\tilde{M} = (\tilde{q}, \tilde{u})$ s.t.

$$\tilde{q}(\theta) := \min\{q^*(\theta), q(\theta)\}$$

$$\text{and } \tilde{u}(\theta) = \int_{\theta}^{\bar{\theta}} \tilde{q}(y) dy$$

- Clearly,
 - \tilde{M} is IC and IR
 - Higher payoff under \tilde{M} than M : \tilde{q} closer to q^{BM} which maximizes virtual surplus
- $\tilde{M} \in \mathcal{M}^{\text{SL}}$?

Sub-optimality of upward adjustments for low costs

- $\tilde{M} \in \mathcal{M}^{\text{SL}} \Leftrightarrow$

$$\underline{W}(\theta, \tilde{q}) := \underline{V}(\tilde{q}(\theta)) - \underline{\theta}\tilde{q}(\theta) - \int_{\theta}^{\bar{\theta}} \tilde{q}(y) dy \geq G^* \quad \forall \theta$$

- Clearly so when
 - $\tilde{q}(\theta) = q(\theta)$ — smaller rents, same TS
 - $\underline{P}^{-1}(\theta) \leq \tilde{q}(\theta) < q(\theta)$ — smaller rents, higher TS
[recall $\underline{P}^{-1}(\theta)$ maximizes $\underline{V}(q) - \underline{\theta}q$]
- $\tilde{q}(\theta) < \min\{q(\theta), \underline{P}^{-1}(\theta)\}$: **not clear**

Sub-optimality of upward adjustments for low costs

- Because $q^*(\bar{\theta}) = q_\ell = \underline{P}^{-1}(\bar{\theta})$, there exists $\theta' > \theta$ s.t.

$$\tilde{q}(y) := \min\{q(y), q^*(\theta)\} \leq \underline{P}^{-1}(y) \quad \forall y \in [\theta, \theta']$$

$$\tilde{q}(\theta') = \min\{\underline{P}^{-1}(\theta'), q(\theta')\}$$

- Monotonicity lemma $\Rightarrow \underline{W}(\cdot, \tilde{q})$ non-increasing over $[\theta, \theta']$
- Previous lemma $\Rightarrow \underline{W}(\theta', \tilde{q}) \geq G^*$
- Hence $\underline{W}(\theta, \tilde{q}) \geq G^*$ Q.E.D.

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Ongoing Work

- **Local robustness**
 - arbitrary set of distributions \mathcal{F}
- Effects of **changes in optimist/pessimism**
 - variations in conjecture (F^* , V^*)
 - variation in demand lower bound \underline{V} (equivalently, \underline{P})
- **Characterization of optimal schedule** when $q^{\text{OPT}}(\theta) < q^{\text{BM}}(\theta)$
 - novel constraint

$$\underline{V}(q(\theta)) - \theta q(\theta) - (1 - \alpha) \int_{\theta}^{\bar{\theta}} q(y) dy \geq G^*$$

- **Comparison with other robustness criteria**

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Conclusions

Conclusions

- **Novel approach to robustness**
 - government has conjecture but does not trust it
 - first protects itself against worst-case
 - then uses conjecture to optimize over worst-case optimal set
- When only uncertainty is over cost
 - robustly optimal mechanism is Baron Myerson with floor
 - efficiency at top and bottom
- Robustness
 - upward quantity adjustment for high cost
 - downward output adjustment for low/intermediate costs
- ...more to be done!

THANK YOU