Money as a Unit of Account

Matthias Doepke and Martin Schneider

Question

- Explain emergence of a common unit of account for future payments.
 - Why coordinate on a common unit of account?
 - What should be the unit of account?

Examples

Treasury Debt, 2002: U.S. Dollars



Examples

Mühlenerbzins, 1794: Meissnische Gulden, bushels of bran



History

- Unit of account often different from medium of exchange.
- Accounting currencies:
 - Distinct from any existing medium of exchange.
 - Livre tournois in France, ECU in Europe.
- Common unit of account in areas with intensive trade:
 - Many currencies used for payment, contracts mostly in one.
 - Vereinsthaler in Northern Germany before unification.
 - Use of dollar denominated contracts in world trade.
- Government-issued fiat money as unit of account:
 - More common recently as governments borrow more . . .
 - ... but not when value too uncertain (dollarization).

Why Coordinate?

- Candidates for unit of account:
 - ► Goods or assets with quoted prices.
- Three features lead to dominant unit of account:
 - 1. Cost of breaking promises.
 - Demand for unit of account that hedges relative-price risk.
 - 2. Trade along credit chains.
 - Demand for common unit of account in chains.
 - 3. Sequential formation of trading networks.
 - Demand for dominant unit of account in entire economy.

What Should Be the Unit of Account?

- Properties of dominant unit of account:
 - Stable in value relative to revenue of borrowers in many transactions.
 - If government is large, government debt works well ...
 - ... but only if value of debt is stable.
 - In areas with a lot of trade, common unit of account is useful: "currency areas."

Literature

- Hedging through denomination of (bilateral) contracts:
 - Bohn (1988), Neumeyer (1999), Schneider-Tornell (2004), Burnside-Eichenbaum-Rebelo (2006) ...
- Credit chains:
 - Kiyotaki-Moore (2001), ...
- Coordination on indexation:
 - Cooper (1990), Acemoglu (1995).
- Medium of exchange and unit of account:
 - Freeman-Tabellini (1998).
- Matching and currency areas:
 - Matsuyama-Kiyotaki-Matsui (1993), Trejos-Wright (2001), Rey (2003) . . .
- Redistribution effects of inflation:
 - Bohn (1990), Doepke-Schneider (2006), Auclert (2006), Doepke et al. (2017) ...

Outline

General setup.

Large default cost and divisible projects:

- Noncontingent contracts, no default, inefficient production.
- Unit of account maximizes scale of production.
- Application to government IOUs.
- Application to optimal currency areas.

Small default cost and indivisible projects (not today):

- Contingent contracts, costly default, efficient production.
- Unit of account minimizes default costs.

Model: Agents, Dates, Goods

• Continuum of agents: Farmers and artisans.

- Meet and write contracts at date 0.
- Work at date 1.
- Exchange goods and consume at date 2.
- Goods:
 - Farm goods: Traded in spot markets at date 2.
 - Artisanal goods: Tailored to matched customer.
 - Labor.

Model: Preferences

Utility function:

$$u_i(c,x,h) = u(\mathbf{c}) + (1+\lambda)x - h.$$

 $u(\mathbf{c})$: Homogeneous utility derived from vector of farm goods \mathbf{c} .

x: Customized artisanal good.

 $h \leq 1$: labor supply.

Model: Farm Endowments and Farm-Goods Market

- Farmer type $i \in \{A, B\}$ with mass 0.5 each.
- ► Farmer of type *i* endowed with one unit of farm good *i* at date 2.
- Farm good i trades in spot market at date 2 at price p_i.
- Price risk: Price of farm good *i* is random.
- ► Vector of farm-good prices p ∈ P is only source of aggregate risk.
- Prices and units of measurement normalized such that utility is linear in wealth and E(p_i) = 1.

Model: Artisan Technology

- ► Mass one each of artisans at location i ∈ [1, 2, ..., N] along highway.
- One unit of labor at date 1 makes one unit of customized artisanal good at date 2.
- Artisans of type 1 produce for farmers, artisans of type *i* + 1 produce for artisans of type *i*.
- Artisanal good valuable only for matched customer.
- Artisanal goods do not trade in spot market and do not have a quoted market price.

Farmers and artisans linked in chains along the highway:

 $\mathsf{Farmer} \longleftarrow \ 1 \longleftarrow \ 2 \longleftarrow \ 3 \longleftarrow \ 4 \longleftarrow \ \ldots \longleftarrow \ \mathsf{N} \ .$

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- Chains created at date 0 by random matching:
 - Morning: Odd *i* artisans travel east and contract with supplier.

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- Chains created at date 0 by random matching:
 - Morning: Odd *i* artisans travel east and contract with supplier.
 - ▶ Night: Odd *i* artisans travel west and contract with customer.

Farmers and artisans linked in chains along the highway:

 $\mathsf{Farmer} \to \leftarrow 1 \to \leftarrow 2 \to \leftarrow 3 \to \leftarrow 4 \to \leftarrow \ldots \to \leftarrow \mathsf{N} \ .$

- Chains created at date 0 by random matching:
 - Morning: Odd i artisans travel east and contract with supplier.
 - ▶ Night: Odd *i* artisans travel west and contract with customer.
 - Matching risk: Identity of farmer in chain unknown in morning.

Model: Contracts

- In every meeting, customer and supplier can enter into contract specifying:
 - 1. Quantity x = h to be produced by supplier in period 1 and delivered in period 2.
 - 2. Payment from customer to supplier in spot market in period 2.
- Introduce friction that favors simple (non-contingent) payment promise:
 - Contract consists of both non-contingent promise and (possibly lower) contingent actual payment.
 - Settling cost if actual payment is lower than promise.
 - Today: Settling cost is infinite: non-contingent promise.

Model: Contracts

- Promise of payment $\pi_{i,j}$:
 - Fixed, non-contingent vector of farm goods.
- Unit of account: Denomination of the promise.

$$oldsymbol{\pi}_{i,j} = oldsymbol{q}_{i,j} \left(egin{array}{c} u_{i,j} \ 1 - u_{i,j} \end{array}
ight).$$

Planning Approach

► To define equilibrium would need to pin down:

- Bargaining process.
- Expectations over contracts in other matches.
- Instead, adopt planning approach:
 - Find system of contracts that maximizes total welfare.
 - Planner chooses (among other things) unit(s) of account for promises.
 - Social optimum is an equilibrium for a specific distribution of bargaining power.

Planning Problem

- Maximizing equally weighted welfare is equivalent to maximizing production of artisanal goods.
- Maximization subject to payment feasibility of payments:
 - If i is artisan with costumer g and supplier j, for any p:

$$\mathbf{p}' \pi_{g,i} \geq \mathbf{p}' \pi_{i,j}.$$

▶ If *i* is farmer with supplier *j*, for any **p**:

$$p_i \geq \mathbf{p}' \boldsymbol{\pi}_{i,j}.$$

Maximization also subject to participation constraints.

Examples for Setup with Large Default Cost

- Assumption on farm good prices:
 - Symmetric price distribution.
 - ► Lower bound of relative price <u>p</u> = min {p_i/p_{-i}} < 1 independent of *i*.
 - Upper bound of relative price $\overline{p} = 1/p > 1$ independent of *i*.

One Farmer, One Artisan: Customized Unit of Account

• One type of farmer and one type of artisan:

$$A \leftarrow 1$$
.

- One stage of matching. Price risk only.
- Decide on $x_A = h_1$ and $\pi_{A,1} = q_{A,1}(u_A, 1 u_A)'$.
- Constraints:
 - Payment feasibility: for all $\mathbf{p} \in \mathbf{P}$,

$$p_a \geq \mathbf{p}' \pi_{A,1}$$

Participation constraints:

$$egin{aligned} 1-q_{\mathcal{A},1}+(1+\lambda)x_{\mathcal{A}}\geq 1,\ q_{\mathcal{A},1}-x_{\mathcal{A}}\geq 0. \end{aligned}$$

One Farmer, One Artisan: Customized Unit of Account

Can achieve first-best production:

- Set artisanal production to $x_A = 1$.
- Make promise in terms of the farmer's good: $u_{A,1} = 1$.
- Scale q_A of payment then has to satisfy:

$$egin{aligned} & p_A \geq p_A q_{A,1}, \ 1-q_{A,1}+1+\lambda \geq 1, \ & q_{A,1}-1 \geq 0. \end{aligned}$$

• Hence, $q_{A,1} = 1$.

 Could not get first-best production with other unit of account. One Farmer, Two Artisans: Unit of Account Passed On

• One type of farmer and two types of artisans:

$$A \longleftarrow 1 \longleftarrow 2.$$

- Two stages of matching. Price risk only.
- Can still achieve first best:
 - Set $x_A = x_1 = 1$.
 - Set $u_{A,1} = u_{1,2} = 1$.
 - Scales of payments need to satisfy:

$$q_{A,1} = q_{1,2} = 1.$$

Two Farmers, Two Artisans: Dominant Unit of Account

Highway with two types of farmer and two types of artisan:

$$\left(\begin{array}{c}A\\B\end{array}\right)\longleftarrow 1\longleftarrow 2.$$

- Two stages of matching. Both price and matching risk.
- Problem: In morning matches of 1 and 2, always possible that night partner of 1 (A or B) will not correspond to the chosen unit of account.
- Scale of production needs to be lowered to avoid default.

Two Farmers, Two Artisans: Dominant Unit of Account

► Consider optimal choice of unit of account *u*, where:

$$\pi_{1,2}=q_{1,2} \left(\begin{array}{c} u\\ 1-u \end{array}\right)$$

The optimal u solves:

$$u = \underset{u}{\operatorname{argmax}} \left\{ \min_{\mathbf{p}} \left\{ \frac{p_i}{p_A u + p_B(1-u)} \right\} \right\}.$$

Under symmetric price distribution have:

$$\min_{\mathbf{p}}\left\{\frac{p_i}{p_A u + p_B(1-u)}\right\} = \frac{\underline{p}}{\max\{u, 1-u\} + \underline{p}\min\{u, 1-u\}},$$

Thus, optimal unit is u = 0.5: Equally weighted bundle of farm goods. Two Farmers, Four Artisans: Dominant Unit of Account

Highway with two types of farmer and four types of artisan:

$$\left(\begin{array}{c}A\\B\end{array}\right)\longleftarrow 1\longleftarrow 2\longleftarrow 3\longleftarrow 4.$$

- Optimal to use equally weighted bundle (u = 0.5) in 3-4 morning matches as well.
- Without coordination on dominant unit of account, additional sources of mismatch, resulting in lower scale of production.

Extensions

- Income risk for farmers: place more weight on good with higher income risk.
- Price distribution not symmetric: farm goods with less volatile prices are better unit of account.
- Small default costs: use unit of account to minimize probability of default.

Decentralization

- Optimal allocation can be decentralized with Nash bargaining at each stage.
- Unit of account is independent of bargaining weights.
- Bargaining weights matter for distribution of surplus across farmers and artisans.

- Model shows that dominant unit of account is optimal.
- In reality, why is money often used, as opposed to a commodity bundle?
- Introduce government that issues IOUs.
- Will private contracts be denominated in government IOUs?

- In period 0, government buys fraction g of farmers' output in exchange for g units of government IOUs.
- ► IOU is claim on tax revenue *T*. Tax revenue is realized at end of date 2, after spot market closes, but before consumption takes place.
- At start of period 2, news about T arrives. IOUs trade in spot market at price:

$$p_{IOU}=E_2(T).$$

- Assume symmetric distribution for p_A and p_B .
- p_{IOU} symmetric with respect to p_A and p_B .
- At extremes of the relative price distribution,

$$rac{p_{IOU}}{\max\{p_A, p_B\}} \in [\underline{p}_{IOU}, \overline{p}_{IOU}], \ \underline{p}_{IOU} < rac{\underline{p}+1}{2}.$$

IOUs can serve as unit of account:

$$\boldsymbol{\pi}_{i,j} = \begin{pmatrix} \pi_{i,j}^{IOU} \\ \pi_{i,j}^{A} \\ \pi_{i,j}^{B} \end{pmatrix} = \boldsymbol{q}_{i,j} \begin{pmatrix} u_{i,j}^{IOU} \\ u_{i,j}^{A} \\ u_{i,j}^{B} \end{pmatrix}$$

Optimal unit of account:

- If $\overline{p}_{IOU} < \frac{\underline{p}+1}{2}$, choose IOUs: $u^{IOU} = 1$.
- Else, choose:

$$u^{IOU} = rac{g}{g + (1 - g)rac{2p}{p+1}},$$

 $u^{A}_{i,j} = u^{B}_{i,j} = rac{1 - u^{IOU}_{i,j}}{2}.$

 Interpretation: "dollarization" when inflation becomes too volatile.

Optimal Currency Areas

Consider model in which there are two locations/countries:

Country A:
$$A \longleftarrow 1 \longleftarrow 2 \longleftarrow 3 \longleftarrow 4$$

Country B: $B \longleftarrow 1 \longleftarrow 2 \longleftarrow 3 \longleftarrow 4$

- At each stage of matching, probability x < 0.5 of meeting someone from the other country.
- ▶ If matched in "wrong" country, can pay cost to rematch.
- Should a common unit of account be adopted?

Optimal Currency Areas

- Separate units (A for A, B for B):
 - Maximizes production conditional on matching within one country.
 - But requires paying rematch cost to avoid possibility of default.
- Common unit of account:
 - Some ex-post risk due to meeting partners from either country.
 - But no need to pay rematch cost.
- Common unit optimal when x sufficiently large.
- Common unit more attractive when chains of credit are longer.

Summary

Three features lead to common unit of account:

- 1. Cost of breaking promises.
- 2. Trade along credit chains.
- 3. Sequential formation of trading networks.
- Properties of optimal unit of account:
 - Stable in value relative to revenue of borrowers in many transactions.
 - Government debt works well if large and not too volatile.
 - Common "currency areas" optimal if lots of trade.

Next Steps

- Explain history of units of accounts and currency areas.
- Examine role of financial intermediaries.
- Examine costs of monetary instability.

Setup with Small Default Cost

- Discrete labor supply $h \in \{0, 1\}$.
- Small default costs: $\kappa < \lambda$.
- Everyone works under optimal allocation.
- Maximize surplus by minimizing probability of default.
- Do this by coordinating on a dominant unit of account.
- Intuition as in large-default-cost case, but rather than extremes of price distribution, probability of default matters.

Optimal Contract

- All agents work: $h_i = 1$ for all *i*.
- Farmers promise and pay their entire harvest.
- Choose promise π in matches between artisans to maximize:

$$E\left[\Pr\left[p_h\left(1+\lambda\right)\geq\mathbf{p}'\boldsymbol{\pi}
ight]
ight]$$

subject to:

$$E\left[\min\left\{p_{h}\left(1+\lambda\right),\mathbf{p}'\boldsymbol{\pi}\right\}\right]\geq1.$$

Actual payment by artisan *i* in chain headed by farmer *h*:

$$m{v}_{i,j}(\mathcal{N},\mathbf{p}) = \min\left\{\mathbf{p}'m{\pi},m{p}_h\left(1+\lambda
ight)
ight\}.$$